Abstract

This paper presents an analysis of entries in the first TAC Market Design Competition final that compares the entries across several scenarios. The analysis complements previous work analyzing the 2007 competition, demonstrating some vulnerabilities of entries that placed highly in the competition. The paper also suggests a simple strategy that would have performed well.

1. Introduction

Auctions are widely used in solving real-world resource allocation problems, and in structuring stock and futures exchanges. As a result, the field of auction mechanism design has drawn much attention in recent years from economists, mathematicians, and computer scientists [2, 10]. In traditional auction theory, auctions are viewed as games of incomplete information, and traditional analytic methods from game theory have been successfully applied to some single-sided auctions, where a single seller has goods for sale and multiple buyers bid for the goods, and some simple forms of double-sided auctions (DAS), where there are multiple sellers and multiple buyers and both sides may make offers or shouts.

However, as, for example Friedman [6], has pointed out, DAS, and particularly continuous double auctions (CDAs),\(^1\) are too complex to analyze in this way since at every moment, a trader must compute expected utility-maximizing shouts based on the history of shouts and transactions, and the time remaining in the auction. This difficulty led researchers to seek experimental approaches. Smith [18] pioneered this field and showed, through a series of experiments with human subjects, that even CDAs with just a handful of traders can give high efficiency and quick convergence to the theoretical equilibrium. Software agents armed with various learning algorithms and optimization techniques have been shown to produce outcomes similar to those obtained by human subjects [4, 9] and are capable of generating higher individual profits [5].

In parallel with the automation of traders, computer scientists have started to take the approach of automated auction mechanism design. Thus, Cliff [3] used a genetic algorithm to discover a completely new kind of auction with good convergence properties. Phelps et al. [15] showed that genetic programming can be used to find an optimal point in a space of pricing policies. Niu et al. [13] identified a mechanism that minimizes variation in transaction price, confirming the mechanism through an evolutionary exploration. Pardoe and Stone [14] suggested a self-adapting auction mechanism that adjusts auction parameters in response to past auction results.

All this work has one common theme — it all studies single markets or compares different market mechanisms indirectly. In contrast, not only do traders in an auction compete against each other, real market institutions face competition [17]. In addition, previous studies usually present comparisons of auction mechanisms in different proprietary settings which differ in available information, computational resources and so on. As a result, mechanisms are difficult to compare, and it is desirable to have some platform that allows multiple markets to compete against each other, and evaluates market mechanisms in a uniform way.

The Trading Agent Competition (TAC) Market Design Tournament [7], also known as the CAT game, addresses these very issues\(^2\). Prior TAC competitions have competing trading agents that aim to maximize their payoffs by interacting in a single market. CAT games do just the opposite. Each entrant in the competition provides a specialist that regulates a market with a set of auction rules, and these specialists compete against each other to attract traders and make profits. Traders in CAT games are provided by the competition platform and each of them learns to choose the best market to trade in.

\(^1\)A CDA is a continuous DA in which any trader can accept an offer and make a deal at any time during the auction period.

\(^2\)The first competition was held in July 2007 and a second competition was held in July 2008.
In a previous paper, we analyzed those entries from the first CAT competition (CAT 2007) that are available in the TAC agent repository, and tried to identify effective auction rules. In this paper, we extend this previous work by more closely examining the relative strength and weakness of the specialist agents.

2. The CAT competition

CAT games are designed to allow markets to compete against each other in a direct fashion. Each market serves a set of traders, each of which makes shouts indicating either what it is prepared to pay to buy some good (a bid), or what it expects to be paid to sell such a good (an ask). A CAT game lasts a certain number of days, each day consists of rounds, and each round lasts a certain number of ticks, or milliseconds.

Each trader is assigned private values for the goods it will trade. For buyers, the private value is the most it will pay for a good. For sellers, the private value is the least it will accept for a good. The private values and the number of goods to buy or sell make up the demand and supply of the markets. Private values remain constant during a day, but may change from day to day, depending upon the configuration of the game server. Each trading agent is endowed with a trading strategy and a market selection strategy. The first specifies how to make offers, the second specifies which market to choose to make shouts in.

Specialists facilitate trade by matching shouts and determining the trading price in an exchange market. Each specialist operates its own exchange market and may choose its own auction rules — the aim of the CAT competition is to create a specialist that optimizes a particular set of measures, including market share, profit share, and transaction success rate. The specialist with the highest cumulative score — the sum of these three metrics — is the winner of a game. Specialists may have adaptive strategies such that the policies change during the course of a game in response to market conditions.

We developed JCAT, the platform that is used in the CAT competition.

2.1. Parameterized market mechanisms

An auction mechanism can be parameterized into components that each regulates an aspect of the market. The following gives a framework extending that in [22]:

- **Matching** policies define the set of matching offers in a market at a given time.
- **Quoting** policies determine the ask quote and bid quote — indicators of where traders need to place asks and bid in order to trade — from existing asks and bids.
- **Shout accepting** policies judge whether a request by a trader to place a shout in the market should be accepted or rejected.
- **Clearing** conditions define when to clear the market and execute transactions between matched shouts.
- **Pricing** policies are responsible for determining transaction prices for matched shouts.
- **Charging** policies determine the charges that a specialist imposes on traders using the market. In JCAT, a specialist can set fees for registration with a specialist, for market information, for making an offer, for completing a transaction, and impose a tax on profits.

For example, the classic CDA mechanism is a combination of the following auction rules (without considering the charging component): the market is cleared whenever a new shout is placed; the market matches the highest bid with the lowest ask that it exceeds; the pricing policy picks the midpoint of a matched ask-bid pair; the quoting policy uses the lowest unmatched ask as the ask quote and the highest unmatched bid as the bid quote; and the shout accepting policy only accepts shouts that beat the corresponding market quote.

2.2. A simple, but powerful market design

We developed MetroCat, a market mechanism that instantiates this parameterized framework, based on the following insights about the CAT game.

- It is crucial to maintain a high transaction success rate, since this rate is not immediately affected by the performance of other markets in contrast to market share and profit share. Thus a strong shout accepting policy, which only allows those shouts that are likely to match with other shouts, is desirable.
- Registration and information fees should be avoided, for these fees cause losses to extra-marginal traders and drive them away. Keeping extra-marginal traders in the market allows them to contribute through their impact on market share.

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4. The number of traders attracted to the market relative to the total number of traders.
5. The amount of profit made by a specialist relative to the total amount of profit made by all the specialists in a game.
6. The number of shouts matched by the market relative to the total number of shouts placed in the market.
7. Traders that theory says should not trade at market equilibrium and will not trade in efficient markets.
• Moderate charges on shouts, transactions, and trader profit only impacts intra-marginal traders, and because of this they stay with the market as long as they can make a considerable amount of profit through transactions after covering fees.

These insights led us to develop a CDA-based market mechanism, which uses a history-based shout accepting policy, denoted as AH. AH is based on the GD trading strategy [8]. GD selects a price that maximizes expected payoff, assuming that, for a given ask price \( a \):

- if another ask price \( a' < a \) was offered and was not accepted by a seller, \( a \) would not be accepted either;
- if another ask price \( a' > a \) was offered and accepted by a seller, \( a \) would have been accepted as well; and
- if a bid price \( b > a \) was offered in the market, \( a \) would have been accepted.

Based on these assumptions, the probability of \( a \) being matched is calculated as:

\[
\Pr(a) = \frac{\sum_{d \geq a} MA(d) + \sum_{d \geq a} B(d)}{\sum_{d \geq a} MA(d) + \sum_{d \geq a} B(d) + \sum_{d \leq a} RA(d)}
\]

where: \( MA(d) \) is the number of asks with price \( d \) that have been matched; \( RA(d) \) is the number of asks with price \( d \) that were not matched, and \( B(d) \) is the number of bids with price \( d \). It is not realistic to keep a full history of shouts and transactions, so GD maintains a sliding window and only considers those shouts and transactions in the window. Computed like this, \( \Pr(a) \) is a monotonically decreasing function, since the higher \( a \) is, the lower \( \Pr(a) \). It is also assumed that when \( a = 0.0 \), \( \Pr(a) = 1 \), and there is a certain value \( u_a \), when \( a > u_a \), \( \Pr(a) = 0 \). The probability \( \Pr(b) \) of a given bid being accepted is computed analogously.

AH uses exactly \( \Pr(a) \) and \( \Pr(b) \) to estimate how likely a shout is to be matched, and only accepts those shouts with a probability higher than a specified threshold \( \lambda \in [0, 1] \). When it is close to 1, the restriction may become too tight for intra-marginal traders to be able to place shouts in the market. When it is close to 0, the restriction may become so loose that extra-marginal traders are able to place shouts which do not stand much chance of being matched. The former would cause both the market and the traders to lose part of the expected profit and lead those traders to leave, and the latter would cause a low transaction success rate. MetroCat uses \( \lambda = 0.5 \), which we found to be optimal for a game configuration similar to CAT 2007.

In addition to AH, MetroCat uses a simple charging policy that imposes low, fixed fees on shouts, transactions, and trader profit, and no charges on registration and information\(^8\). Since we developed the competition platform, MetroCat was not an entrant in the competition\(^9\), but we have used it as a benchmark in our post-competition experiments.

3. The analysis of CAT entries

Trading competitions have been an effective tool in fostering innovative approaches and advocating enthusiasm and exchange among researchers [20, 21]. However, the competitions themselves usually cannot provide a complete view of the relative strength and weakness of entries. In a competition, the performance of one player closely depends upon the composition of its opponents and the competition configuration, and the scenarios considered are usually limited. Thus we typically turn to post-competition analysis to tell us which entries are most interesting.

Ideally, such an analysis will cover all possible scenarios, but this usually presents too large a possible space. As a result, a common practice is to deliberately select a limited number of representative strategies and run games corresponding to a set of discrete points in the infinite space, assuming that the results are representative of what would happen in the whole space were one to explore it [19].

There are two common types of approaches to post-competition analysis: white-box approaches and black-box approaches. A white-box approach attempts to relate the internal logic and features of strategies and game outcomes. A black-box approach considers strategies as atomic entities. Niu et al. [12] performed a white-box analysis of CAT 2007, and examined how the dynamics in the CAT games are affected by the policies of each entry and their adaptation over time. This paper takes a complementary black-box approach and also examines the weakness and strength of CAT 2007 entries against MetroCat.

3.1. Multi-lateral simulation

Now, a full analysis of a set of strategies can only be achieved by considering many runs (to eliminate randomness) of every possible combination of strategies. This is not feasible for the CAT competition where each game runs for around 5 hours (irrespective of the hardware — the length of each trading days is hardcoded at a constant that permits each specialist to take time to perform possibly complex computations — any reduction in this time would potentially distort the results). Inspired by ecological analyses like [1, 16] — in which more copies of successful strategies, and less copies of unsuccessful strategies are run for

\(^8\)The fees MetroCat imposes on shouts, transactions, and trader profit are respectively 0.1, 0.1, and 10% during the post-tournament experiments described in later sections.

\(^9\)Instead it was included in the JCAT source code provided to entrants in the first CAT competition to support the development of their entries.
each successive game — but constrained by the number of specialists that we could have in a single game, we modified each strategy’s playing time in proportion to its score. That is, in a game that included all specialists, we decreased the number of trading days for less successful strategies, and increased the days for more successful strategies.

Figures 1(a) and 1(b) show the result of this simulation. The distribution on the y-axis shows the proportion of the total number of trading days for all markets that are allotted to each market, indicating how this evolves in populations without and with MetroCat respectively. Figure 1(a) shows that without MetroCat (i) the results of this analysis agrees with the results reported in [12], again confirming that IAMwildCAT was the strongest entrant in the 2007 competition; and (ii) the days allotted to PersianCat shrink more slowly than those of other losing specialists. This agrees with the results of bilateral games between IAMwildCAT and PersianCat (described below) and suggests that PersianCat was a strong entry. Figure 1(b) shows that with MetroCat (i) MetroCat quickly dominates the other entries, doing so faster than IAMwildCAT in Figure 1(a), and by generation 8 only MetroCat has any trading days; and (ii) the days allotted to MetroCat have any trading days; and (ii) the CAT 2007 champion, IAMwildCAT, loses trading days faster than those for other specialists. This agrees with the results of bilateral games between IAMwildCAT and PersianCat (described below) and suggests that PersianCat was a strong entry. Figure 1(b) shows that with MetroCat (i) MetroCat quickly dominates the other entries, doing so faster than IAMwildCAT in Figure 1(a), and by generation 8 only MetroCat has any trading days; and (ii) the CAT 2007 champion, IAMwildCAT, loses trading days faster than those for other specialists. This agrees with the results of bilateral games between IAMwildCAT and PersianCat (described below) and suggests that PersianCat was a strong entry.

3.2. Bilateral simulation

One-on-one games closely examine the strength and weakness of a specialist when it faces different opponents. As a result, we ran 81 experiments in total between the 9 specialists including MetroCat. Nine of these are self-play games. Table 1 shows the resulting payoffs of specialists — their average daily scores — in these CAT games. Each payoff is averaged over 10 iterations and entry \((i, j)\) is the payoff of specialist \(i\) in the game against specialist \(j\).

Figure 3 compares these payoffs pictorially using a polar coordinate system. Each plot shows the 9 specialists evenly distributed on the outer circle, the radial coordinates of the 9 vertices of the solid-line polygon represent a given specialist’s payoffs against all 9 specialists, and the radial coordinates of the 9 vertices of the dashed-line polygon represent its opponents’ payoffs in these games. The solid-line polygon and the dashed-line polygon overlap on the vertex that corresponds to the self-play game of the particular specialist. In Figure 3(a), the solid-line polygon completely encloses the dashed-line one, meaning MetroCat wins over all the other 8 specialists in bilateral competitions\(^{10}\). Figure 3(i) shows the opposite situation in that Mertacor loses all the games. The two polygons for any other specialist intersect somehow, showing their advantages in some games and disadvantages in others.

Both Figure 3(b) and 3(h) show that IAMwildCAT, the CAT 2007 champion, surprisingly loses (although barely) against PersianCat which placed sixth. This explains why in Figure 1(a) the days for PersianCat shrink more slowly than those for other specialists. IAMwildCAT’s loss, given the defeat of PersianCat by PSUCAT and jackaroo, suggests that IAMwildCAT has some particular weakness that is taken advantage of by PersianCat.

Other discrepancies, when compared to the results of the 2007 competition, include jackaroo (which placed 4th) winning over PSUCAT (2nd) and Crocodile (3rd). These may be significant, or may be caused by differences in the configurations for PSUCAT and Crocodile used in the sim-

\(^{10}\) MetroCat maintains a better balance than those CAT 2007 entries between market share and profit share by keeping extra-marginal traders and preventing them from placing uncompetitive shouts.
3.3 Ecological simulation

The payoff table for bilateral CAT games can be used to approximate ecological dynamics for populations involving more than 2 specialist types. The payoff of each specialist type for a certain population mixture is computed as the expected payoff for this specialist assuming that each specialist obtains the payoff it would have obtained had it competed one-on-one with each of the other specialists in the mix. Under this assumption, Figure 2(a) and 2(b) show how a population with an initial even distribution of specialists evolves over time when, as in [1], every specialist plays against every other specialist in every generation in bilateral games, and the number of specialists in any generation is proportional to the payoff achieved by that “breed” of specialist in the previous generation.

Comparing Figure 2(a) with Figure 1(a) and Figure 2(b) with Figure 1(b), shows that while the winning strategies are the same, the ecological simulations based on multi-lateral games converge much faster than those based on bilateral games. This may be explained by the more epidemic effects of the strength of particular specialists in multi-lateral games compared with bilateral games. Another noticeable phenomenon is that PSUCAT performs much worse in the simulations with bilateral games than those with multi-lateral games, while jackaroo and IAMwildCAT do the opposite. These discrepancies indicate that, as one might expect, different game setups may lead to very different results. However, our results may be helpful to identify the weakness in strategies by looking at the particular scenarios in which a strategy performs poorly.

4. Conclusions and future work

This paper reports results of post-competition simulations of entries in the First TAC Market Design Competition...
based on both bilateral games and multi-lateral games. The results basically agree with those in [12] and the result of the actual tournament, and also unveil weaknesses of specialists in particular scenarios, including the defeat of IAMwildCAT by PersianCat in bilateral games and the poorer relative performance of jackaroo in multi-lateral scenarios than in bilateral scenarios.

Some simulations also consider an additional specialist, MetroCat, which uses a history-based shout accepting policy that is derived from the GD trading strategy for double auctions. MetroCat claims victories in all the scenarios addressed in this paper, showing the importance of a shout accepting policy in a market mechanism.

The bilateral games and multi-lateral games can be viewed as the two ends of a spectrum of CAT games. The aim of running simulations based on both configurations is to explore whether the different competition configurations lead to results that differ very much. It is hoped that if they make no much difference — and our results suggest that they do not — the low cost of bilateral games can be used to approximate the games involving more different individual types and different population distributions.

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References

Figure 3: Payoffs of self and opponents in bilateral CAT games. On the outer circles starting from polar angle $0^\circ$ lists the 9 specialists anti-clockwise: PersianCat ($0^\circ$), MANX ($40^\circ$), jackaroo ($80^\circ$), IAMwildCAT ($120^\circ$), MetroCat ($160^\circ$), PSUCAT ($200^\circ$), CrocodileAgent ($240^\circ$), TacTex ($280^\circ$), and Mertacor ($320^\circ$). The radial coordinates of the 9 vertices of a solid-line polygon represent the corresponding specialist’s payoffs against the 9 specialists respectively, and those of a dashed-line polygon represent payoffs of its opponents respectively. The overlapping vertex of the two polygons in each sub-figure is the self-play game for the particular specialist.