RECURSION

Recursion

- Recursion is defining something in terms of itself
- There are many examples in nature:
 - Seeds in a sunflower





- ...in mathematics:
 - Factorial
 - Induction
- ... and in computer graphics:
 - Koch snowflake

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Today

- This lecture looks at
 - The basics of recursion.
 - Some examples of recursive functions.
- The textbook doesn't cover recursion in any detail (the only material is on pages 96 and 97 in my copy)..

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Koch snowflake

- Starting with a line, then:
 - 1. Divide each line into three segments of equal length.
 - 2. Draw an equilateral triangle that has the middle segment from step 1 as its base and points outward.
 - 3. Remove the line segment that is the base of the triangle from step 2.
- Repeat as often as you like.

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• Here are the first four iterations of the Koch snowflake.









• The more iterations, the more snowflaky it looks.

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Here it is in C++

```
// r1.cpp
#include <iostream>
using namespace std;

int power( int x, int y ) {
  if ( y == 0 )
    return( 1 );
  else
    return( x * power( x, y-1 ));
} // end of power()

int main() {
  cout << "2^3 = " << power( 2,3 ) << endl;
}</pre>
```

Power function

• *Power* is defined recursively:

$$x^{y} = \begin{cases} \text{if } y == 0, \ x^{y} = 1 \\ \text{otherwise, } x^{y} = x * x^{y-1} \end{cases}$$

- There are two parts to the definition:
 - The *base case*, what we do when *y* is zero.
 - The *recursive case*, what we do when *y* is not zero.
- This is the common pattern for all recursive definitions.

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- Notice that power() calls itself!
- This seems to be magic, but we'll see how it is done in a moment.
- You can make recursive calls with any method *except main()*
- BUT beware of infinite loops!!!
- You have to know when and how to stop the recursion what is the *stopping* condition.

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Walking through power (2,4)

• Initial call is power (2, 4)

	call	x	y	return value
1	power(2,4)	2	4	2 * power(2,3)
2	power(2,3)	2	3	2 * power(2,2)
3	power(2,2)	2	2	2 * power(2,1)
4	power(2,1)	2	1	2 * power(2, 0)
4	power(2,0)	2	0	1

- The first is the *original call*
- Followed by four *recursive calls*

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Another example: factorial

• *factorial* is defined recursively:

$$N! = \begin{cases} \text{if } N == 1, & N! = 1\\ \text{otherwise}, & N! = N * (N-1)! \end{cases}$$
(for $N > 0$)

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Stacks

- The computer uses a data structure called a *stack* to keep track of what is going on
- Think of a *stack* like a stack of plates
- You can only take off the top one
- You can only add more plates to the top
- This corresponds to the two basic *stack operations*:
 - *push* putting something onto the stack
 - pop taking something off of the stack
- When each recursive call is made, power() is pushed onto the stack
- When each return is made, the corresponding power () is popped off of the stack

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Here it is in C++

```
// r2.cpp
#include <iostream>
using namespace std;

int factorial ( int N ) {
   if ( N == 1 )
      return( 1 );
   else
      return( N * factorial( N-1 ));
} // end of factorial()

int main() {
   cout << "5! = " << factorial( 5 ) << endl;
}</pre>
```

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• Walk through factorial(4)

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• What is the output of this program?

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Another example

```
//r3.cpp
#include <iostream>
using namespace std;

void countDown (int n) {
   if ( n <= 0 )
      cout << "Blastoff!" << endl;
   else {
      cout << "Time to launch is " << n << " seconds" << endl;
      countDown(n - 1);
   }
} // end of countDown()

int main() {
   countDown(5);
}

cisl5-fall2007-parsons-lectVI.1</pre>
```

• Now, let's switch the statements in the recursive case around.

```
//r4.cpp
#include <iostream>
using namespace std;

void countDown (int n) {
   if ( n <= 0 )
      cout << "Blastoff!" << endl;
   else {
      countDown(n - 1);
      cout << "Time to launch is " << n << " seconds" << endl;
   }
} // end of countDown()

int main() {
   countDown(5);
}

cisl5-fall2007-parsons-lectVI.1</pre>
```

• What is the output of this program?

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Fibonacci

```
// in r5.cpp
int fibonacci (int n) {
   if (n == 0){
     return 0;
   }
   else
     if (n == 1){
       return 1;
    }
    else {
      return(fibonacci(n - 1) + fibonacci(n -2));
   }
} // end of fibonacci()
```

• Again countDown has the general structure:

```
// base case part
if (<base-case condition>)
   return <base-case-value>
// general case
```

else

return <recursively computed expression>

• This is common to all recursive functions — the only difference you'll see is that some functions have two base cases.

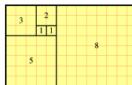
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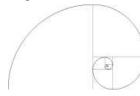
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• A tiling where tile sides are successive members of the Fibonacci sequence.



- A spiral constructed from the above tiling.



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Recursive and iteration

- You can use recursion to iterate.
- (Iteration = repetition, what you would normally do with a loop).
- The following slide has an example.

void array::printI() {

cout << endl;

} // end of printI()

 $if (x < size) {$

printR(x+1);

cout << endl;</pre>

} // end of printR()

else {

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for (int i=0; i<size; i++)</pre>

cout << data[i] << " ";</pre>

void array::printR(int x) {

cout << data[x] << " ";</pre>

• Compare printI() (iterative) with printR() (recursive).

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```

```
// recursion-iteration.cpp
   #include <iostream>
   using namespace std;
   class array {
   private:
     int *data;
     int size;
   public:
     array( int n ) : size( n ) { data = new int[n]; }
     void set( int x, int v ) { data[x] = v; }
     void printI();
     void printR( int x );
   }; // end of class array
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```

```
int main() {
     array A( 5 );
     for ( int i=0; i<5; i++ )
       A.set( i,i*10 );
     cout << "output from iterative printI(): ";</pre>
     A.printI();
     cout << "output from recursive printR(): ";</pre>
     A.printR( 0 );
   } // end of main() method
   and the output is:
   output from iterative printI(): 0 10 20 30 40
   output from recursive printR(): 0 10 20 30 40
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```

And the details...

• In the recursive version, each call is like one iteration inside the for loop in the iterative version

	call	index	output	next call
1	printR(0)	0	0	printR(1)
2	printR(1)	1	10	printR(2)
	printR(2)	2	20	printR(3)
4	printR(3)	3	30	printR(4)
5	printR(4)	4	40	printR(5)
6	printR(5)	5	endl	(none)

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Summary

- This lecture has looked at
 - The basic idea of recursion
 - A bunch of different examples of recursion
- You will find that the idea of recursion gets easier to cope with as you get more familiar with it.
- That means, like all programming ideas, it gets easier with use.

- With recursion, each time the function is invoked, one step is taken towards the resolution of the task the function is meant to complete.
- Before each step is executed, the state of the task being completed is somewhere in the middle of being completed.
- After each step, the state of the task is one step closer to completion.
- In the example above, each time printR(i) is called, the array is printed from the *i*-th element to the end of the array.
- In the power (x,y) example, each time the function is called, power is computed for each x^y , in terms of the previous x^{y-1} .
- In the factorial (N) example, each time the function is called, factorial is computed for each N, in terms of the previous N-1.

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