## PROBLEM SOLVING AGENTS

## **Problem Solving Agents**

- Lecture 1 introduced *rational agents*.
- Now consider agents as *problem solvers*:
- Systems which set themselves *goals* and find *sequences of actions* that achieve these goals.
- What is a problem? A *goal* and a *means* for achieving the goal.
- The goal specifies the state of affairs we want to bring about.
- The means specifies the operations we can perform in an attempt to bring about the means.
- The difficulty is deciding which operations and what *order* to carry out the operations.

```
cis32-spring2009-parsons-lect07
```

# Overview

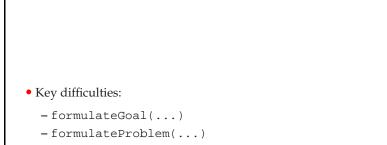
#### Aims of the this lecture:

- introduce problem solving;
- introduce *goal formulation*;
- show how problems can be stated as *state space search*;
- show the importance and role of *abstraction*;
- introduce *undirected search*:
  - breadth 1st search;
  - depth 1st search.
- define main performance measures for search.

```
cis32-spring2009-parsons-lect07
```

# • Operation of problem solving agent: /\* s is sequence of actions \*/ repeat { percept = observeWorld(); state = updateState(state, p); if s is empty then { goal = formulateGoal(state); prob = formulateProblem(state, goal); s = search(prob); } action = first(s); s = remainder(s); } until false; /\* i.e., forever \*/

cis32-spring2009-parsons-lect07



- $-\operatorname{search}(\ldots)$
- It isn't easy to see how to tackle any of these.
- Here we will concentrate mainly on search.

#### cis32-spring2009-parsons-lect07

# **Problem Formulation**

- Once goal is determined, formulate the problem to be solved.
- First determine set of possible states *S* of the problem.
- Then problem has:
  - *initial state —* the starting point, *s*<sub>0</sub>;
  - *operations* the actions that can be performed,  $\{o_1, \ldots, o_n\}$ .
  - *goal* what you are aiming at subset of *S*.

# **Goal Formulation**

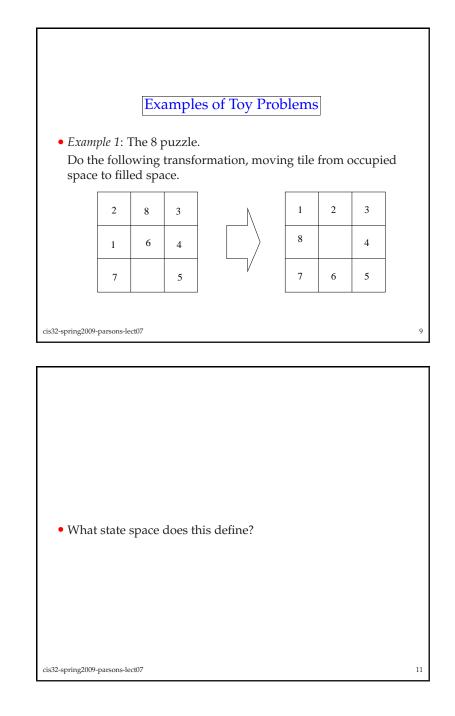
- Where do an agent's goals come from?
  - Agent is a *program* with a *specification*.
  - Specification is to maximise performance measure.
  - Should *adopt goal* if achievement of that goal will maximise this measure.
- Goals provide a *focus* and *filter* for decision-making:
  - *focus*: need to consider how to achieve them;
  - *filter*: need not consider actions that are incompatible with goals.
- For this course, we will assume that an agent is given its goals.

cis32-spring2009-parsons-lect07

- The initial state together with operations determines *state space* of problem.
- Operations cause *changes* in state.
- Solution is a sequence of actions such that when applied to initial state *s*<sub>0</sub>, we have goal state.
- What does this look like?

cis32-spring2009-parsons-lect07

7



- Initial state as shown above.
- Goal state as shown above.
- Operations:
  - $o_1$ : move any tile to left of empty square to right;
  - $-o_2$ : ?
  - $-o_3$ : ?
- $-o_4$ : ?

cis32-spring2009-parsons-lect07

- Example 2: The *n* queens problem from chess.
- Place *n* queens on chess board so that no queen can be taken by another.
- Initial state: empty chess board.
- Goal state: *n* queens on chess board, one occupying each space, so that none can take others.
- Operations: place queen in empty square.

		-
Solution Cost		• As an examp
• For most problems, some solutions are better than others:		
<ul> <li>in 8 puzzle, number of moves to get to solution;</li> <li>number of moves to checkmate;</li> <li>length of distance to travel.</li> </ul>		
• Mechanism for determining <i>cost</i> of solution is <i>path cost function</i>		
• This is the length of the path through the state-space from the initial state to the goal state.		• How many n
is32-spring2009-parsons-lect07	13	cis32-spring2009-parsons-lect0
• There are five moves:		
1. 2. 3.		<ul> <li>In the state sp a path throug</li> <li>When we sold of what const</li> </ul>
4.		• It is hard to p
5.		• Instead we st
• What are they?		computers ar
• What does the path through the solution space look like?		• A <i>brute force</i> a <i>searching</i> thro one that achi
cis32-spring2009-parsons-lect07	15	cis32-spring2009-parsons-lect0

ble, consider the following state in the 8-puzzle:

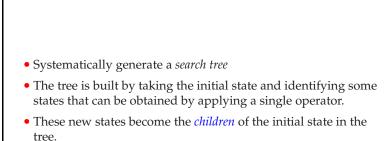
2	8	3
1	6	4
7		5

noves are there to the solution?

Problem Solving as Search

- pace view of the world, finding a solution is finding gh the state space.
- lve a problem like the 8-puzzle we have some idea stitutes the next best move.
- program this kind of approach.
- tart by programming the kind of repetitive task that re good at.
- approach to problem solving involves *exhaustively* ough the space of *all possible* action sequences to find ieves goal.

s32-spring2009-parsons-lect07



- These new states are then examined to see if they are the goal state.
- If not, the process is repeated on the new states.
- We can formalise this description by giving an algorithm for it.

```
cis32-spring2009-parsons-lect07
```

• Note the difference between *state space* and *search tree*.

- State space is every possible state and the relationships between them.
  - It is static.
- Search tree the set of states the agent has looked at (is looking at) and some of the relationships between them.
  - It is dynamic.

### • General algorithm for search:

agenda = initial state; while agenda not empty do{ pick node from agenda; new nodes = apply operations to state; if goal state in new nodes then { return solution; } add new nodes to agenda; }

18

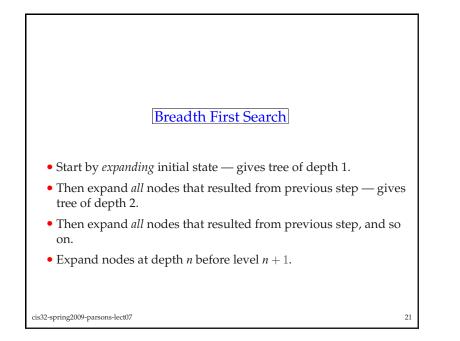
20

cis32-spring2009-parsons-lect07

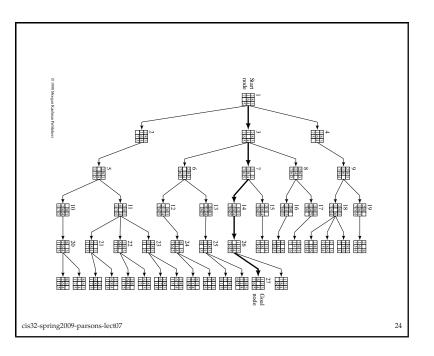
17

19

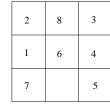
- Question: How to pick states for expansion?
- Two obvious solutions:
- depth first search;
- breadth first search.



```
/* Breadth first search */
agenda = initial state;
while agenda not empty do
{
    pick node from front of agenda;
    new nodes = apply operations to state;
    if goal state in new nodes then
    {
        return solution;
    }
    APPEND new nodes to END of agenda;
}
```



• For the 8-puzzle as so:



23

• We have the following state space:

• Given this numbering of the states, the agenda would look like
1. 1
2. 2, 3, 4
3. 3, 4, 5
4. 4, 5, 6, 7, 8
5. 5, 6, 7, 8, 9
6. 6, 7, 8, 9, 10, 11.
7
cis32-spring2009-parsons-lect07

- Advantage: *guaranteed* to reach a solution if one exists.
- If all solutions occur at depth *n*, then this is good approach.
- Disadvantage: time taken to reach solution!
- Let *b* be *branching factor* average number of operations that may be performed from any level.
- If solution occurs at depth *d*, then we will look at

 $1+b+b^2+\cdots+b^d$ 

nodes before reaching solution — *exponential*.

cis32-spring2009-parsons-lect07

25

27

#### • Time for breadth first search:

Depth	Nodes	Time
0	1	1 msec
1	11	.01 sec
2	111	.1 sec
4	11,111	11 secs
6	$10^{6}$	18 mins
8	$10^{8}$	31 hours
10	$10^{10}$	128 days
12	$10^{12}$	35 years
14	$10^{14}$	2500 years
20	$10^{20}$	$3^{15}$ years

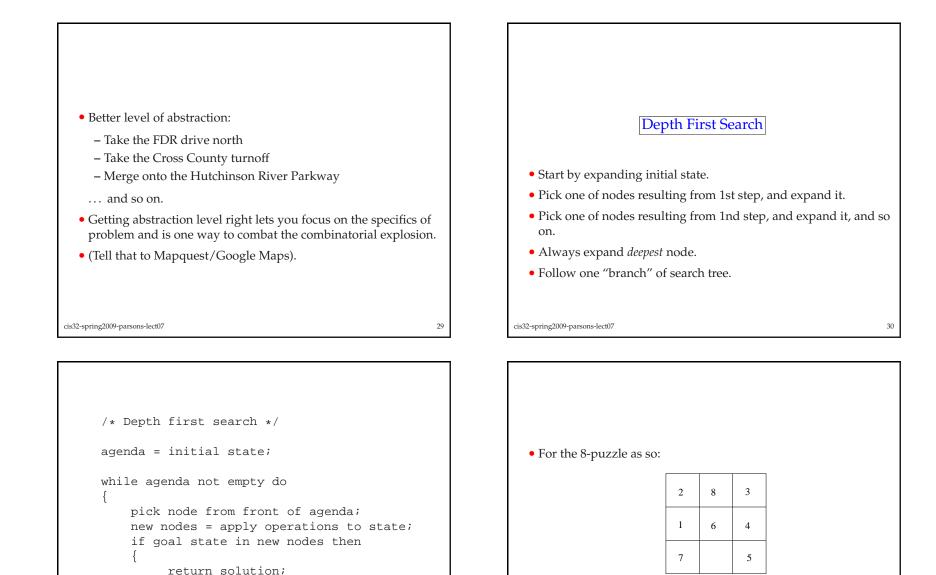
• *Combinatorial explosion!* 

cis32-spring2009-parsons-lect07

# Importance of ABSTRACTION

- When formulating a problem, it is crucial to pick the right level of *abstraction*.
- Example: Given the task of driving from where I used live in Manhatatan to Boston.
- Some possible actions...
  - depress clutch;
  - turn steering wheel right 10 degrees;
- ... inappropriate level of *abstraction*. Too much *irrelevant detail*.

cis32-spring2009-parsons-lect07

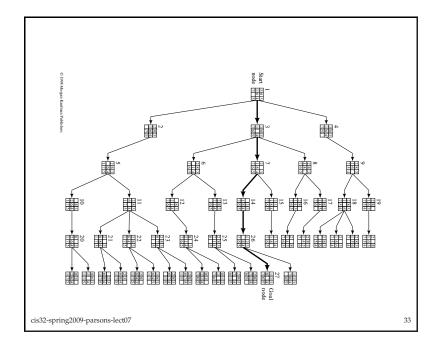


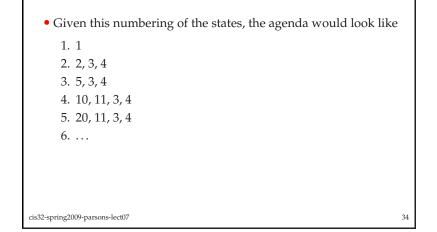
31

put new nodes on FRONT of agenda;

cis32-spring2009-parsons-lect07

• We have the following state space:





Performance Measures for Search

• Completeness:

Is the search technique *guaranteed* to find a solution if one exists?

- *Time complexity*: How many computations are required to find solution?
- *Space complexity:* How much memory space is required?
- *Optimality*: How good is a solution going to be w.r.t. the path cost function.

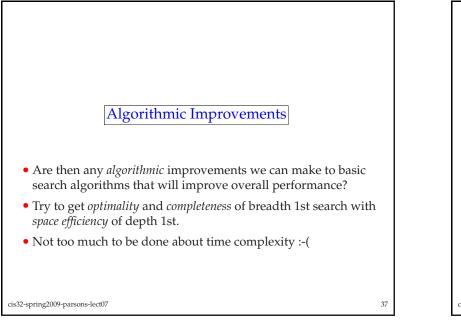
36

cis32-spring2009-parsons-lect07

35

• Depth first search is *not* guaranteed to find a solution if one exists.

- However, if it *does* find one, amount of time taken is much less than breadth first search.
- *Memory requirement* is much less than breadth first search.
- Solution found is *not* guaranteed to be the best.



# Depth Limited Search

- Depth first search has some desirable properties space complexity.
- But if wrong branch is expanded (with no solution on it), then it won't terminate.
- Idea: introduce a *depth limit* on branches to be expanded.
- Don't expand a branch below this depth.

```
cis32-spring2009-parsons-lect07
```

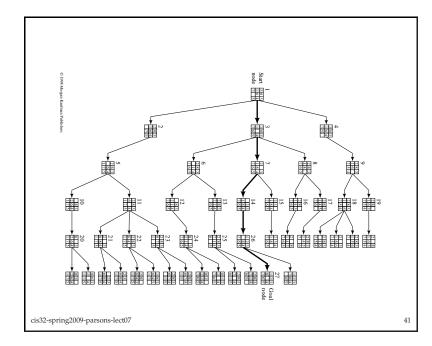
• General algorithm for depth limited search: depth limit = max depth to search to; agenda = initial state; while agenda not empty do take node from front of agenda; new nodes = apply operations to node; if goal state in new nodes then { return solution; } if depth(node) < depth limit then { add new nodes to front of agenda; } }

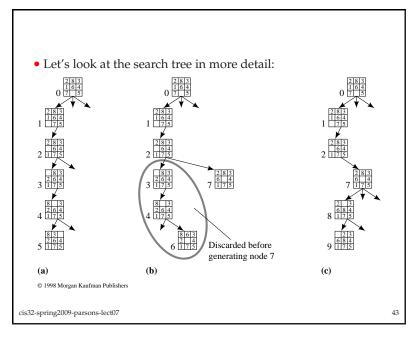
#### • For the 8-puzzle as so:

2	8	3
1	6	4
7		5

• We have the following state space:

cis32-spring2009-parsons-lect07





• Given this numbering of the states, a depth limited search with depth limit of three would have an agenda that looks like

1
 2, 3, 4
 5, 3, 4
 4, 10, 11, 3, 4
 5, 11, 3, 4
 6, 3, 4
 7, 6, 7, 8, 4
 8, 12, 13, 7, 8, 4
 9, 13, 7, 8, 4
 10....

cis32-spring2009-parsons-lect07

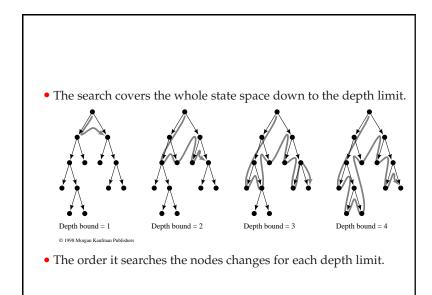
- So, when we hit the depth bound, we don't add any more nodes to the agenda.
- Then we pick the next node off the agenda.
- This has the effect of moving the search back to the last node above depth limit that is "partly expanded".
- This is known as *chronological backtracking*.
- The effect of the depth limit is to force the search of the whole state space down to the limit.
- We get the completeness of breadth-first (down to the limit), with the space cost of depth first.

cis32-spring2009-parsons-lect07

# **Iterative Deepening**

- Unfortunately, if we choose a max depth for d.l.s. such that shortest solution is longer, d.l.s. is not complete.
- Iterative deepening an ingenious *complete* version of it.
- Basic idea is:
  - do d.l.s. for depth 1; if solution found, return it;
  - otherwise do d.l.s. for depth n; if solution found, return it; otherwise, ...
- So we *repeat* d.l.s. for all depths until solution found.

```
cis32-spring2009-parsons-lect07
```



```
General algorithm for iterative deepening search:
depth limit = 1;
repeat {
    result = depth_limited_search(
        max depth = depth limit;
        agenda = initial node;
    );
    if result contains goal then {
        return result;
    }
    depth limit = depth limit + 1;
    } until false; /* i.e., forever */
Calls d.l.s. as subroutine.
```

```
cis32-spring2009-parsons-lect07
```

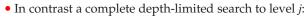
45

47

- Note that in iterative deepening, we *re-generate nodes on the fly*.
   Each time we do call on depth limited search for depth *d*, we need to regenerate the tree to depth *d* − 1.
- Isn't this inefficient?
- Tradeoff *time* for *memory*.
- In general we might take a *little* more time, but we save a *lot* of memory.
- Now for breadth-first search to level *d*:

$$N_{bf} = 1 + b + b^2 + \dots b^d$$
  
=  $rac{b^{d+1} - 1}{b - 1}$ 

cis32-spring2009-parsons-lect07



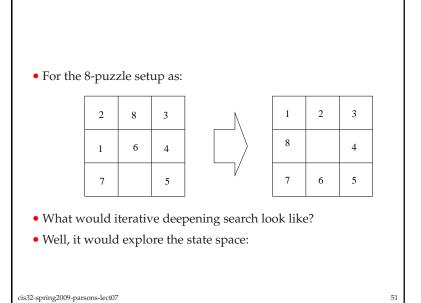
$$N_{df}^{j} = rac{b^{j+1}-1}{b-1}$$

• (This is just a breadth-first search to depth *j*.)

• In the worst case, then we have to do this to depth *d*, so expanding:

$$N_{id} = \sum_{j=0}^{d} \frac{b^{j+1} - 1}{b - 1}$$
  
:  
$$= \frac{b^{d+2} - 2b - bd + d + 1}{(b - 1)^2}$$

cis32-spring2009-parsons-lect07



• For large *d*:

$$\frac{N_{id}}{N_{bf}} = \frac{b}{b-1}$$

- So for high branching and relatively deep goals we do a small amount more work.
- Example: Suppose b = 10 and d = 5.

Breadth first search would require examining 111, 111 nodes, with memory requirement of 100, 000 nodes.

Iterative deepening for same problem: 123, 456 nodes to be searched, with memory requirement only 50 nodes. Takes 11% longer in this case.

50

cis32-spring2009-parsons-lect07

