PROBLEM SOLVING AGENTS

Overview

Aims of the this lecture:

- introduce *problem solving*;
- introduce *goal formulation*;
- show how problems can be stated as *state space search*;
- show the importance and role of *abstraction*;
- introduce *undirected search*:
 - breadth 1st search;
 - depth 1st search.
- define main performance measures for search.

Problem Solving Agents

- Lecture 1 introduced *rational agents*.
- Now consider agents as *problem solvers*: Systems which set themselves *goals* and find *sequences of actions* that achieve these goals.
- What is a problem?

A *goal* and a *means* for achieving the goal.

- The goal specifies the state of affairs we want to bring about.
- The means specifies the operations we can perform in an attempt to bring about the means.
- The difficulty is deciding which operations and what *order* to carry out the operations.

```
• Operation of problem solving agent:
/* s is sequence of actions */
repeat {
    percept = observeWorld();
    state = updateState(state, p);
    if s is empty then {
         goal = formulateGoal(state);
         prob = formulateProblem(state, goal);
         s = search(prob);
    action = first(s);
    s = remainder(s);
until false; /* i.e., forever */
```

- Key difficulties:
 - formulateGoal(...)
 - formulateProblem(...)
 - **-**search(...)
- It isn't easy to see how to tackle any of these.
- Here we will concentrate mainly on search.

Goal Formulation

- Where do an agent's goals come from?
 - Agent is a *program* with a *specification*.
 - Specification is to maximise performance measure.
 - Should *adopt goal* if achievement of that goal will maximise this measure.
- Goals provide a *focus* and *filter* for decision-making:
 - *focus*: need to consider how to achieve them;
 - *filter*: need not consider actions that are incompatible with goals.
- For this course, we will assume that an agent is given its goals.

Problem Formulation

- Once goal is determined, formulate the problem to be solved.
- First determine set of possible states *S* of the problem.
- Then problem has:
 - *initial state —* the starting point, *s*₀;
 - *operations* the actions that can be performed, $\{o_1, \ldots, o_n\}$.
 - *goal* what you are aiming at subset of *S*.

- The initial state together with operations determines *state space* of problem.
- Operations cause *changes* in state.
- Solution is a sequence of actions such that when applied to initial state *s*₀, we have goal state.
- What does this look like?

Examples of Toy Problems

• *Example 1*: The 8 puzzle.

Do the following transformation, moving tile from occupied space to filled space.



- Initial state as shown above.
- Goal state as shown above.
- Operations:
 - o_1 : move any tile to left of empty square to right;

$$-o_2$$
: ?

- $-o_3$: ?
- $-o_4$: ?

• What state space does this define?

- Example 2: The *n* queens problem from chess.
- Place *n* queens on chess board so that no queen can be taken by another.
- Initial state: empty chess board.
- Goal state: *n* queens on chess board, one occupying each space, so that none can take others.
- Operations: place queen in empty square.

Solution Cost

- For most problems, some solutions are better than others:
 - in 8 puzzle, number of moves to get to solution;
 - number of moves to checkmate;
 - length of distance to travel.
- Mechanism for determining *cost* of solution is *path cost function*.
- This is the length of the path through the state-space from the initial state to the goal state.

• As an example, consider the following state in the 8-puzzle:

2	8	3
1	6	4
7		5

• How many moves are there to the solution?

- There are five moves:
 - 1.
 - 2.
 - 3.
 - **J**.
 - 4.
 - 5.
- What are they?
- What does the path through the solution space look like?

Problem Solving as Search

- In the state space view of the world, finding a solution is finding a path through the state space.
- When we solve a problem like the 8-puzzle we have some idea of what constitutes the next best move.
- It is hard to program this kind of approach.
- Instead we start by programming the kind of repetitive task that computers are good at.
- A *brute force* approach to problem solving involves *exhaustively searching* through the space of *all possible* action sequences to find one that achieves goal.

- Systematically generate a *search tree*
- The tree is built by taking the initial state and identifying some states that can be obtained by applying a single operator.
- These new states become the *children* of the initial state in the tree.
- These new states are then examined to see if they are the goal state.
- If not, the process is repeated on the new states.
- We can formalise this description by giving an algorithm for it.

```
• General algorithm for search:
```

```
agenda = initial state;
while agenda not empty do{
    pick node from agenda;
    new nodes = apply operations to state;
    if goal state in new nodes
    then {
        return solution;
     }
    add new nodes to agenda;
}
```

- Note the difference between *state space* and *search tree*.
- State space is every possible state and the relationships between them.
 - It is static.
- Search tree the set of states the agent has looked at (is looking at) and some of the relationships between them.
 - It is dynamic.

- Question: How to pick states for expansion?
- Two obvious solutions:
 - depth first search;
 - breadth first search.

Breadth First Search

- Start by *expanding* initial state gives tree of depth 1.
- Then expand *all* nodes that resulted from previous step gives tree of depth 2.
- Then expand *all* nodes that resulted from previous step, and so on.
- Expand nodes at depth *n* before level n + 1.

```
/* Breadth first search */
   agenda = initial state;
   while agenda not empty do
       pick node from front of agenda;
       new nodes = apply operations to state;
       if goal state in new nodes then
             return solution;
       APPEND new nodes to END of agenda;
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```

• For the 8-puzzle as so:

2	8	3
1	6	4
7		5

• We have the following state space:



• Given this numbering of the states, the agenda would look like

1. 1
 2. 2, 3, 4
 3. 3, 4, 5
 4. 4, 5, 6, 7, 8
 5. 5, 6, 7, 8, 9
 6. 6, 7, 8, 9, 10, 11.
 7. ...

- Advantage: *guaranteed* to reach a solution if one exists.
- If all solutions occur at depth *n*, then this is good approach.
- Disadvantage: time taken to reach solution!
- Let *b* be *branching factor* average number of operations that may be performed from any level.
- If solution occurs at depth *d*, then we will look at

 $1+b+b^2+\cdots+b^d$

nodes before reaching solution — *exponential*.

• Time for breadth first search:

Depth	Nodes	Time
0	1	1 msec
1	11	.01 sec
2	111	.1 sec
4	11,111	11 secs
6	10^{6}	18 mins
8	10^{8}	31 hours
10	10^{10}	128 days
12	10^{12}	35 years
14	10^{14}	2500 years
20	10^{20}	3^{15} years
	1	

• *Combinatorial explosion!*

Importance of ABSTRACTION

- When formulating a problem, it is crucial to pick the right level of *abstraction*.
- Example: Given the task of driving from where I used live in Manhatatan to Boston.
- Some possible actions...
 - depress clutch;
 - turn steering wheel right 10 degrees;
 - ... inappropriate level of *abstraction*. Too much *irrelevant detail*.

- Better level of abstraction:
 - Take the FDR drive north
 - Take the Cross County turnoff
 - Merge onto the Hutchinson River Parkway
 - ... and so on.
- Getting abstraction level right lets you focus on the specifics of problem and is one way to combat the combinatorial explosion.
- (Tell that to Mapquest/Google Maps).

Depth First Search

- Start by expanding initial state.
- Pick one of nodes resulting from 1st step, and expand it.
- Pick one of nodes resulting from 1nd step, and expand it, and so on.
- Always expand *deepest* node.
- Follow one "branch" of search tree.

```
/* Depth first search */
agenda = initial state;
while agenda not empty do
    pick node from front of agenda;
    new nodes = apply operations to state;
    if goal state in new nodes then
         return solution;
put new nodes on FRONT of agenda;
```

• For the 8-puzzle as so:

2	8	3
1	6	4
7		5

• We have the following state space:



• Given this numbering of the states, the agenda would look like

1
 2, 3, 4
 5, 3, 4
 10, 11, 3, 4
 20, 11, 3, 4

6. ...

- Depth first search is *not* guaranteed to find a solution if one exists.
- However, if it *does* find one, amount of time taken is much less than breadth first search.
- *Memory requirement* is much less than breadth first search.
- Solution found is *not* guaranteed to be the best.

Performance Measures for Search

• Completeness:

Is the search technique *guaranteed* to find a solution if one exists?

• *Time complexity*:

How many computations are required to find solution?

• *Space complexity*:

How much memory space is required?

• *Optimality*:

How good is a solution going to be w.r.t. the path cost function.
Algorithmic Improvements

- Are then any *algorithmic* improvements we can make to basic search algorithms that will improve overall performance?
- Try to get *optimality* and *completeness* of breadth 1st search with *space efficiency* of depth 1st.
- Not too much to be done about time complexity :-(

Depth Limited Search

- Depth first search has some desirable properties space complexity.
- But if wrong branch is expanded (with no solution on it), then it won't terminate.
- Idea: introduce a *depth limit* on branches to be expanded.
- Don't expand a branch below this depth.

```
• General algorithm for depth limited search:
```

```
depth limit = max depth to search to;
agenda = initial state;
while agenda not empty do
  take node from front of agenda;
  new nodes = apply operations to node;
  if goal state in new nodes then {
    return solution;
  if depth(node) < depth limit then {
    add new nodes to front of agenda;
```

• For the 8-puzzle as so:

2	8	3
1	6	4
7		5

• We have the following state space:



• Given this numbering of the states, a depth limited search with depth limit of three would have an agenda that looks like

1. 1
 2. 2, 3, 4
 3. 5, 3, 4
 4. 10, 11, 3, 4
 5. 11, 3, 4
 6. 3, 4
 7. 6, 7, 8, 4
 8. 12, 13, 7, 8, 4
 9. 13, 7, 8, 4
 10....







- So, when we hit the depth bound, we don't add any more nodes to the agenda.
- Then we pick the next node off the agenda.
- This has the effect of moving the search back to the last node above depth limit that that is "partly expanded".
- This is known as *chronological backtracking*.
- The effect of the depth limit is to force the search of the whole state space down to the limit.
- We get the completeness of breadth-first (down to the limit), with the space cost of depth first.

Iterative Deepening

- Unfortunately, if we choose a max depth for d.l.s. such that shortest solution is longer, d.l.s. is not complete.
- Iterative deepening an ingenious *complete* version of it.
- Basic idea is:
 - do d.l.s. for depth 1; if solution found, return it;
 - otherwise do d.l.s. for depth n; if solution found, return it;
 - otherwise, ...
- So we *repeat* d.l.s. for all depths until solution found.

• General algorithm for iterative deepening search:

```
depth limit = 1;
repeat {
  result = depth_limited_search(
    max depth = depth limit;
    agenda = initial node;
  );
  if result contains goal then {
    return result;
  }
  depth limit = depth limit + 1;
} until false; /* i.e., forever */
```

• Calls d.l.s. as subroutine.

• The search covers the whole state space down to the depth limit.



• The order it searches the nodes changes for each depth limit.

- Note that in iterative deepening, we *re-generate nodes on the fly*.
 Each time we do call on depth limited search for depth *d*, we need to regenerate the tree to depth *d* − 1.
- Isn't this inefficient?
- Tradeoff *time* for *memory*.
- In general we might take a *little* more time, but we save a *lot* of memory.
- Now for breadth-first search to level *d*:

$$N_{bf} = 1 + b + b^2 + \dots b^d$$

= $\frac{b^{d+1} - 1}{b - 1}$

• In contrast a complete depth-limited search to level *j*:

$$N_{df}^{j} = \frac{b^{j+1} - 1}{b - 1}$$

- (This is just a breadth-first search to depth *j*.)
- In the worst case, then we have to do this to depth *d*, so expanding:

$$V_{id} = \sum_{j=0}^{d} \frac{b^{j+1} - 1}{b - 1}$$

= $\frac{b^{d+2} - 2b - bd + d + 1}{(b - 1)^2}$

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• For large *d*:

$$\frac{N_{id}}{N_{bf}} = \frac{b}{b-1}$$

- So for high branching and relatively deep goals we do a small amount more work.
- Example: Suppose b = 10 and d = 5.

Breadth first search would require examining 111, 111 nodes, with memory requirement of 100, 000 nodes.

Iterative deepening for same problem: 123, 456 nodes to be searched, with memory requirement only 50 nodes.

Takes 11% longer in this case.

• For the 8-puzzle setup as:



- What would iterative deepening search look like?
- Well, it would explore the state space:



• In the following way.

• States would be expanded in the order:

- 1. 1
 2. 1, 2, 3, 4
 3. 1, 2, 5, 3, 6, 7, 8, 4, 9.
 4. 1, 2, 5, 10, 11, 3, 6, 12, 13, 7, 14, 15, 8, 16, 17, 4, 9, 18, 19.
 5. ...
- Note that these are the states *visited*, not the nodes on the agenda.

Bi-directional Search

- Suppose we search from *the goal state backwards* as well as from *initial state forwards*.
- Involves determining *predecessor* nodes to goal, and then looking at predecessor nodes to this, ...
- Rather than doing one search of b^d , we do *two* $b^{d/2}$ searches.
- Much more efficient.

• Example:

Suppose b = 10, d = 6.

Breadth first search will examine Bidirectional search will examine

nodes.

- Can combine different search strategies in different directions.
- For large *d*, is still impractical!

Summary

- This lecture introduced the basics of problem solving.
- In particular it discussed *state space* models and looked at the basic techniques for solving them.
 - Search for the goal.
 - Path through state space is the solution.
- We also looked at two techniques for search:
 - Breadth first.
 - Depth first.