

- ... which interact through communication ...
- ... are able to act in an environment ...
- ... have different "spheres of influence" (which may coincide) ...
- ... will be linked by other (organisational) relationships.

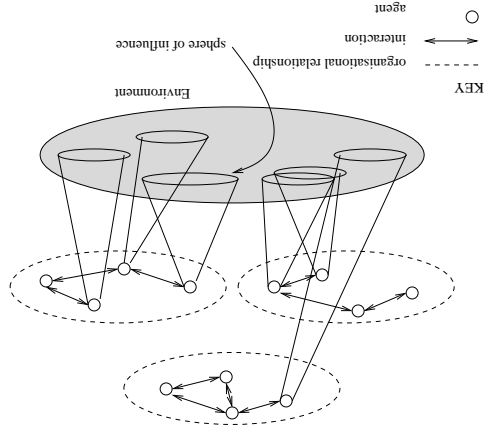
LECTURE 6: MULTIAGENT INTERACTIONS

An Introduction to Multiagent Systems

- Assume we have just two agents: $Ag = \{i, j\}$.
- Agents are assumed to be *self-interested*: they *have preferences over how the environment is*.
- Assume $\Omega = \{\omega_1, \omega_2, \dots\}$ is the set of "outcomes" that agents have preferences over.
- We capture preferences by *utility functions*:
 $u_i: \Omega \rightarrow \mathbb{R}$
 $u_j: \Omega \rightarrow \mathbb{R}$
- Utility functions lead to *preference orderings* over outcomes:

$\omega \succ_i \omega'$ means $u_i(\omega) > u_i(\omega')$
 $\omega \succeq_i \omega'$ means $u_i(\omega) \geq u_i(\omega')$

2 Utilities and Preferences



1 What are Multiagent Systems?

3 Multiagent Encounters

- We need a model of the environment in which these agents will act...
- agents simultaneously choose an action to perform, and as a result of the actions they select, an outcome in Ω will result;
- the *actual* outcome depends on the *combination* of actions;
- assume each agent has just two possible actions that it can perform C (“cooperate”) and “D” (“defect”).
- Environment behaviour given by *state transformer function*:

$$\tau : \hat{A}_C \times \hat{A}_C \rightarrow \Omega$$

agent's action agent's action

Rational Action

- Suppose we have the case where *both* agents can influence the outcome, and they have utility functions as follows:
- With a bit of abuse of notation:

$$u_i(\omega_1) = 1 \quad u_i(\omega_2) = 1 \quad u_i(\omega_3) = 4 \quad u_i(\omega_4) = 4$$

$$u_j(\omega_1) = 1 \quad u_j(\omega_2) = 4 \quad u_j(\omega_3) = 1 \quad u_j(\omega_4) = 4$$

$$u_i(D, D) = 1 \quad u_i(D, C) = 1 \quad u_i(C, D) = 4 \quad u_i(C, C) = 4$$

$$u_j(D, D) = 1 \quad u_j(D, C) = 4 \quad u_j(C, D) = 1 \quad u_j(C, C) = 4$$

- Then agent *i*'s preferences are:

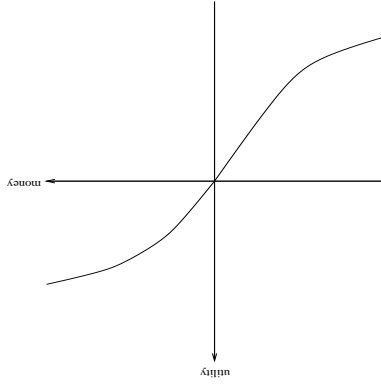
$$C, C \preceq_i C, D \succ_i D, C \preceq_i D, D$$

- “*C*” is the *rational choice* for *i*.

(Because *i* prefers all outcomes that arise through *C* over all outcomes that arise through *D*.)

What is Utility?

- Utility is *not* money (but it is a useful analogy).
- Typical relationship between utility & money:



- Here is a state transformer function:

$$\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_3 \quad \tau(C, C) = \omega_4$$

- (This environment is sensitive to actions of both agents.)

- Here is another:

$$\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_1 \quad \tau(C, D) = \omega_1 \quad \tau(C, C) = \omega_1$$

- (Neither agent has any influence in this environment.)

- And here is another:

$$\tau(D, D) = \omega_1 \quad \tau(D, C) = \omega_2 \quad \tau(C, D) = \omega_1 \quad \tau(C, C) = \omega_2$$

- (This environment is controlled by *j*.)

Dominant Strategies

- Given any particular strategy s (either C or D) agent i , there will be a number of possible outcomes.
- We say s_1 *dominates* s_2 if every outcome possible by i playing s_1 is preferred over every outcome possible by i playing s_2 .
- A rational agent will never play a dominated strategy.
- So in deciding what to do, we can *delete dominated strategies*.
- Unfortunately, there isn't always a unique undominated strategy.

Payoff Matrices

- We can characterise the previous scenario in a *payoff matrix*

	i	
	defect	coop
j	defect	1 1
	coop	1 4

- Agent i is the *column player*.
- Agent j is the *row player*.

Nash Equilibrium

- In general, we will say that two strategies s_1 and s_2 are in Nash equilibrium if:
 - under the assumption that agent i plays s_1 , agent j can do no better than play s_2 ; and
 - under the assumption that agent j plays s_2 , agent i can do no better than play s_1 .
- Neither agent has any incentive to deviate from a Nash equilibrium.
- Unfortunately:

- Not every interaction scenario has a Nash equilibrium.
- Some interaction scenarios have more than one Nash equilibrium.

Competitive and Zero-Sum Interactions

- Where preferences of agents are diametrically opposed we have *strictly competitive* scenarios.
- Zero-sum encounters are those where utilities sum to zero:

$$u_i(\omega) + u_j(\omega) = 0 \quad \text{for all } \omega \in \Omega.$$
- Zero sum implies strictly competitive.
- Zero sum encounters in real life are very rare . . . but people tend to act in many scenarios as if they were zero sum.

- This apparent paradox is *the fundamental problem of multi-agent interactions*.

It appears to imply that *cooperation will not occur in societies of self-interested agents*.

- Real world examples:

- nuclear arms reduction (“why don’t I keep mine...”)
- free rider systems — public transport;
- in the UK — television licenses.
- The prisoner’s dilemma is *ubiquitous*.
- Can we recover cooperation?

- The *individual rational* action is *defect*.
- This guarantees a payoff of no worse than 2, whereas cooperating guarantees a payoff of at most 1.
- So defection is the best response to all possible strategies: both agents defect, and get payoff = 2.
- But *intuition* says this is *not* the best outcome: Surely they should both cooperate and each get payoff of 3!

- Payoff matrix for prisoner’s dilemma:

	<i>i</i>	
	defect	coop
<i>j</i>	defect	2, 1
	coop	4, 3

- Top left: If both defect, then both get punishment for mutual defection.
- Top right: If *i* cooperates and *j* defects, *i* gets sucker’s payoff of 1, while *j* gets 4.
- Bottom left: If *j* cooperates and *i* defects, *j* gets sucker’s payoff of 1, while *i* gets 4.
- Bottom right: Reward for mutual cooperation.

4 The Prisoner’s Dilemma

Two men are collectively charged with a crime and held in separate cells, with no way of meeting or communicating. They are told that:

- if one confesses and the other does not, the confessor will be freed, and the other will be jailed for three years;
- if both confess, then each will be jailed for two years.

Both prisoners know that if neither confesses, then they will each be jailed for one year.

Arguments for Recovering Cooperation

- Conclusions that some have drawn from this analysis:
 - the game theory notion of rational action is wrong!
 - somehow the dilemma is being formulated wrongly
- Arguments to recover cooperation:
 - We are not all machiavelli!
 - The other prisoner is my twin!
 - The shadow of the future. . .

4.2 Backwards Induction

- But. . . suppose you both know that you will play the game exactly n times.
 - On round $n - 1$, you have an incentive to defect, to gain that extra bit of payoff. . .
 - But this makes round $n - 2$ the last "real", and so you have an incentive to defect there, too.
- This is the *backwards induction* problem.
 - Playing the prisoner's dilemma with a fixed, finite, pre-determined, commonly known number of rounds, defection is the best strategy.

4.1 The Iterated Prisoner's Dilemma

- One answer: *play the game more than once*.
 - If you know you will be meeting your opponent again, then the incentive to defect appears to evaporate.
- *Cooperation is the rational choice in the infinitely repeated prisoner's dilemma.* (Hurrah!)

4.3 Axelrod's Tournament

- Suppose you play iterated prisoner's dilemma against a *range* of opponents. . .
 - What strategy should you choose, so as to maximise your overall payoff?
- Axelrod (1984) investigated this problem, with a computer tournament for programs playing the prisoner's dilemma.

5 Game of Chicken

Consider another type of encounter — the *game of chicken*:

	defect	coop
defect	1, 1	4, 3
coop	2, 2	3, 3

(Think of James Dean in *Rebel without a Cause*: swerving = coop, driving straight = defect.)

Difference to prisoner's dilemma:

Mutual defection is most feared outcome.

(Whereas sucker's payoff is most feared in prisoner's dilemma.)
 Strategies (c,d) and (d,c) are in Nash equilibrium

Strategies in Axelrod's Tournament

- ALLD: "Always defect" — the *hawk* strategy;
- TIT-FOR-TAT:
- 1. On round $n = 0$, cooperate.
- 2. On round $n > 0$, do what your opponent did on round $n - 1$.
- TESTER:
- On 1st round, defect. If the opponent retaliated, then play TIT-FOR-TAT. Otherwise intersperse cooperation & defection.
- JOSS:
- As TIT-FOR-TAT, except periodically defect.

6 Other Symmetric 2 x 2 Games

Given the 4 possible outcomes of (symmetric) cooperate/defect games, there are 24 possible orderings on outcomes.

– CC >! CD >! DC >! DD

Cooperation dominates.

– DC >! DD >! CC >! CD

Deadlock. You will always do best by defecting.

– DC >! CC >! DD >! CD

Prisoner's dilemma.

– DC >! CC >! CD >! DD

Chicken.

– CC >! DC >! DD >! CD

Stag hunt.

Recipes for Success in Axelrod's Tournament

- Axelrod suggests the following rules for succeeding in his tournament:
- *Don't be envious:*
Don't play as if it were zero sum!
- *Be nice:*
Start by cooperating, and reciprocate cooperation.
- *Retaliate appropriately:*
Always punish defection immediately, but use "measured" force — don't overdo it.
- *Don't hold grudges:*
Always reciprocate cooperation immediately.