EVOLUTION AND CO-EVOLUTION OF COMPUTER PROGRAMS TO CONTROL INDEPENDENTLY-ACTING AGENTS

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ABSTRACT

This paper describes the recently developed "genetic programming" paradigm which genetically breeds populations of computer programs to solve problems. In genetic programming, the individuals in the population are hierarchical computer programs of various sizes and shapes. This paper also extends the genetic programming paradigm to a "co-evolution" algorithm which operates simultaneously on two populations of independently-acting hierarchical computer programs of various sizes and shapes.

1. INTRODUCTION AND OVERVIEW

This paper describes the recently developed "genetic programming" paradigm which genetically breeds populations of computer programs to solve problems. In genetic programming, the individuals in the population are hierarchical compositions of functions and arguments of various sizes and shapes. Each of these individual computer programs is evaluated for its fitness in handling the problem environment and a simulated evolutionary process is driven by this measure of fitness.

This paper also extends the genetic programming paradigm to a "co-evolution" algorithm which operates simultaneously on two (or more) populations of independently-acting hierarchical computer programs of various sizes and shapes. In co-evolution, each population acts as the environment for the other population. In particular, each individual of the first population is evaluated for "relative fitness" by testing it against each individual in the second population, and, simultaneously, each individual in the second population is evaluated for "relative fitness" by testing it against each individual in the first population. Over a period of many generations, individuals with high "absolute fitness" tend to evolve as the two populations mutually bootstrap each other to increasingly high levels of fitness.

In this paper, the genetic programming paradigm is illustrated with three problems. The first problem involves genetically breeding a population of computer programs to allow an "artificial ant" to traverse an irregular trail. The second problem involves genetically breeding a minimax control strategy in a differential game with an independently-acting pursuer and evader. The third problem illustrates the "co-evolution" and involves genetically breeding an optimal strategy for a player of a simple discrete two-person game represented by a game tree in extensive form.

2. BACKGROUND ON GENETIC ALGORITHMS

Genetic algorithms are highly parallel mathematical algorithms that transform populations of individual mathematical objects (typically fixed-length binary character strings) into new populations using operations patterned after (1) natural genetic operations such as sexual recombination (crossover) and (2) fitness proportionate reproduction (Darwinian survival of the fittest). Genetic algorithms begin with an initial population of individuals (typically randomly generated) and then iteratively (1) evaluate the individuals in the population for fitness with respect to the problem environment and (2) perform genetic operations on various individuals in the population to produce a new population. John Holland of the University of Michigan presented the pioneering formulation of genetic algorithms for fixed-length character strings in Adaptation in Natural and Artificial Systems (Holland 1975). Holland established, among other things, that the genetic algorithm is a mathematically near optimal approach to adaptation in that it maximizes expected overall average payoff when the adaptive process is viewed as a multi-armed slot machine problem requiring an optimal allocation of future trials given currently available information. Recent work in genetic algorithms and genetic classifier systems can be surveyed in Goldberg (1989), Davis (1987), and Schaffer (1989).

3. BACKGROUND ON GENETIC PROGRAMMING PARADIGM

Representation is a key issue in genetic algorithm work because genetic algorithms directly manipulate the coded representation of the problem and because the representation scheme can severely limit the window by which the system observes its world. Fixed length character strings present difficulties for some problems — particularly problems in artificial intelligence where the desired solution is hierarchical and where the size and shape of the solution is unknown in advance. The need for more powerful representations has been recognized for some time (De Jong 1985, 1988).

The structure of the individual mathematical objects that are manipulated by the genetic algorithm can be more complex than the fixed length character strings. Smith (1980) departed from the early fixed-length character strings by introducing variable length strings, including strings whose elements were if-then rules (rather than single characters). Holland’s introduction of the classifier system (1986) continued the trend towards increasing the complexity of the
structures undergoing adaptation. The classifier system is a cognitive architecture into which the genetic algorithm is embedded so as to allow adaptive modification of a population of string-based if-then rules (whose condition and action parts are fixed length binary strings). Applications of classifier systems are discussed in Wilson (1987a, 1987b, 1988).

In addition, we have recently shown that entire computer programs can be genetically bred to solve problems in a variety of different areas of artificial intelligence, machine learning, and symbolic processing (Koza 1989, 1990a). In this recently developed "genetic programming" paradigm, the individuals in the population are compositions of functions and terminals appropriate to the particular problem domain. The set of functions used typically includes arithmetic operations, mathematical functions, conditional logical operations, and domain-specific functions. Each function in the function set must be well defined for any element in the range of every other function in the set. The set of terminals used typically includes inputs (sensors) appropriate to the problem domain and various constants. The search space is the hyperspace of all possible compositions of functions that can be recursively composed of the available functions and terminals. The symbolic expressions (S-expressions) of the LISP programming language are an especially convenient way to create and manipulate the compositions of functions and terminals described above. These S-expressions in LISP correspond directly to the "parse tree" that is internally created by most compilers.

The basic genetic operations for the genetic programming paradigm are fitness proportionate reproduction and crossover (recombination). Fitness proportionate reproduction is the basic engine of Darwinian reproduction and survival of the fittest and operates for genetic programming paradigms in the same way as it does for conventional genetic algorithms. The crossover operation for genetic programming paradigms is a sexual operation that operates on two parental LISP S-expressions and produces two offspring S-expressions using parts of each parent. In particular, the crossover operation creates new offspring S-expressions by exchanging sub-trees (i.e. sub-lists) between the two parents. Because entire sub-trees are swapped, this genetic crossover (recombination) operation produces syntactically and semantically valid LISP S-expressions as offspring regardless of which point is selected in either parent.

For example, consider the parental LISP S-expression:

\[(\text{OR} \ (\text{NOT} \ D1) \ (\text{AND} \ D0 \ D1))\]

And, consider the second parental S-expression below:

\[(\text{OR} \ D1 \ (\text{NOT} \ D0))\]

\[(\text{AND} \ (\text{NOT} \ D0) \ (\text{NOT} \ D1))\]

These two LISP S-expressions can be depicted graphically as rooted, point-labeled trees with ordered branches. Assume that the points of both trees are numbered in a depth-first way starting at the left. Suppose that the second point (out of 6 points of the first parent) is randomly selected as the crossover point for the first parent and that the sixth point (out of 10 points of the second parent) is randomly selected as the crossover point of the second parent. The crossover points are therefore the NOT in the first parent and the AND in the second parent.

The two parental LISP S-expressions are shown below:

```
(OR
 (NOT D1)
 (AND D0 D1))
```

```
(OR
 D1
 (AND (NOT D0) (NOT D1))
)
```

The two crossover fragments are two sub-trees shown below:

```
NOT
 D1
```

```
AND
 D0
```

These two crossover fragments correspond to the bold, underlined sub-expressions (sub-lists) in the two parental LISP S-expressions shown above. The two offspring resulting from crossover are shown below.

```
(OR
 (AND (NOT D0) (NOT D1))
 (AND D0 D1))
```

```
(OR
 D0
 D1
)
```

Note that the first offspring above is a perfect solution for the exclusive-or function, namely

\[(\text{OR} \ (\text{AND} \ (\text{NOT} \ D0) \ (\text{NOT} \ D1)) \ (\text{AND} \ D0 \ D1))\].

Details on this algorithm can be found in Koza (1990a).

This new genetic algorithm paradigm has been successfully applied (Koza 1989, 1990a) to example problems in several different areas, including (1) machine learning of functions (e.g. learning the Boolean 11-multiplexer function), (2) planning (e.g. developing a robotic action sequence that can stack an arbitrary initial configuration of blocks into a specified order), (3) automatic programming (e.g. discovering a computational procedure for solving pairs of linear equations, solving quadratic equations for complex roots, and discovering trigonometric identities), (4) sequence induction (e.g. inducing a recursive computational procedure for the Fibonacci and the Hofstadter sequences), (5) pattern recognition (e.g. translation-invariant recognition of a simple one-dimensional shape in a linear retina), (6) optimal control (e.g. centering a cart and balancing a broom on a moving cart in minimal time by applying a "bang bang"
force to the cart) (Koza and Keane 1990a, Koza and Keane 1990b), (7) symbolic "data to function" regression, symbolic "data to function" integration, and symbolic "data to function" differentiation, (8) symbolic solution to functional equations (including differential equations with initial conditions, integral equations, and general functional equations), (9) empirical discovery (e.g. rediscovering Kepler's Third Law, rediscovering the well-known econometric "exchange equation" $MV = PQ$ from actual time series data for the money supply, the velocity of money, the price level, and the gross national product of an economy) (Koza 1990b), and (10) simultaneous architectural design and training of neural networks.

4. THE "ARTIFICIAL ANT"

Jefferson et al. (1990) devised a planning task involving an "artificial ant" attempting to traverse an irregular trail and successfully used the conventional string-based genetic algorithm to discover a finite state automaton enabling the "artificial ant" to traverse the trail.

The setting for the problem is a square 32 by 32 toroidal grid in the plane. The "Santa Fe" trail is a winding trail of stones in 89 of the 1024 cells. The trail has single missing stones, double missing stones, single missing stones at some corners, double missing stones at other corners (knight moves), and triple missing stones at other corners (long knight moves).

The "artificial ant" begins in the cell identified by the coordinates (0,0) and is facing in a particular direction (i.e. east). The artificial ant has a sensor that can see only the single adjacent cell in the direction the ant is currently facing. At each time step, the ant has the capacity to execute any of four operations, namely, to move forward (advance) in the direction it is facing, to turn right (and not move), to turn left (and not move), or to sense the contents of the single adjacent cell in the direction the ant is facing.

The objective of the ant is to traverse the entire trail. Jefferson, Collins et al. limited the ant to a certain number of time steps (200).

Jefferson, Collins et al. started by assuming that the finite automaton necessary to solve the problem would have 32 or fewer states. They then represented an individual in their population of automata by a binary string representing the state transition diagram (and its initial state) of the individual automaton. The ant's sensory input at each time step was coded as one bit and the output at each time step was coded as two bits (representing the three actions). The next state of the automaton was coded with 5 bits. The complete behavior of an automaton was thus specified with a genome consisting of a binary string with 453 bits (64 substrings of length 7 representing the state transitions plus 5 additional bits representing the initial state of the automaton). They then processed a population of 65,536 individual bit strings of length 453 on a Connection Machine™ using a genetic algorithm using crossover and mutation operating on a selected (relatively small) fraction of the population. After 200 generations in a particular run (taking about 10 hours on the Connection Machine), Jefferson, Collins et al. reported that a single individual in the population emerged which attained a perfect score of 89 stones. As it happened, this single individual completed the task in exactly 200 operations.

In our approach to this task using the genetic programming paradigm, we used the function set consisting of the functions $F = \{\text{IF-SENSOR, PROGN}\}$. The IF-SENSOR function has two arguments and evaluates the first argument if the ant's sensor senses a stone or, otherwise, evaluates the second argument. The PROGN function is the LISP connective function that sequentially evaluates its arguments as individual steps in a "program." The terminal set was $T = \{\text{ADVANCE, TURN-RIGHT, TURN-LEFT}\}$. These three terminals are actually functions with no arguments. They operate via their side effects on the ant's state (i.e. the ant's position on the grid or the ant's facing direction). IF-SENSOR, ADVANCE, TURN-RIGHT, and TURN-LEFT correspond directly to the operators defined and used by Jefferson, Collins et al. We allowed 400 time steps before timing out.

The initial generation (generation 0) consisted of randomly generated individual S-expressions recursively created using the available functions and available terminals of the problem. For this problem (and each of the other problems described in this paper), each new generation was created from the preceding generation by applying the fitness proportionate reproduction operation to 10% of the population and by applying the crossover operation to 90% of the population (with reselection allowed). The selection of crossover points in the population was biased 90% towards internal (function) points of the tree and 10% towards external (terminal) points of the tree. Mutation was not used. For practical reasons, a limit of 4 was placed on the depth of initial random S-expressions and a limit of 15 was placed on the depth of S-expressions created by crossover.

In one run, an individual LISP S-expression scoring 89 out of 89 emerged on the 7th generation, namely,

\[
\text{IF-SENSOR (ADVANCE)}
\]
This plan is graphically depicted in Figure 1.

This individual LISP S-expression is the solution to the problem. In particular, this plan moves the ant forward if a stone is sensed. Otherwise it turns right and then moves the ant forward if a stone is sensed but turns left (returning to its original orientation) if no stone is sensed. Then it turns left and moves forward if a stone is sensed but turns right (returning to its original orientation) if no stone is sensed. If the ant originally did not sense a stone, the ant moves forward unconditionally as its fifth operation. Note that there is no testing of the backwards directions.

We can undertake to measure the number of individuals that need to be processed by a genetic algorithm to produce a desired result with a certain probability, say 99%. Suppose, for example, a particular run of a genetic algorithm produces the desired result with only a probability of success $p_s$ after a specified choice (perhaps arbitrary and non-optimal) of number of generations $N_{gen}$ and population of size $N$. Suppose also that we are seeking to achieve the desired result with a probability of, say, $z = 1 - \varepsilon = 99%$. Then, the number $K$ of independent runs required is

$$K = \frac{\log(1-z)}{\log(1-p_s)} = \frac{\log \varepsilon}{\log(1-p_s)},$$

where $\varepsilon = 1-z$.

For example, we ran 111 runs of the Artificial Ant problem with a population size of 1000 and 51 generations. We found that the probability of success $p_s$ on a particular single run was 43%. With this probability of success, 8 independent runs are required to assure a 99% probability of solving the problem on at least one of the 8 runs. That is, it is sufficient to process 408,000 individuals. This requires about 6 hours of computing time on the Texas Instruments Explorer II+™ workstation.

The graph below shows that the probability of success $p_s$ of a run is 67% for a population size of 2000 with 51 generations. The graph also shows that the probability of success of a run is 81% for a population of 4000 with 51 generations. A population of 2000 requires 4 independent runs (i.e. 408,000 individuals are sufficient) to achieve the desired 99% probability. A population of 4000 requires 3 independent runs (i.e. 612,000 individuals are sufficient).

**Figure. 1 Artificial Ant-Santa Fe Trail**

### 5. DIFFERENTIAL PURSUIT GAME

As a second illustration of the genetic programming paradigm, we consider a differential pursuer-evader game. Our objective is to find an optimal strategy for one player when the environment (fitness function) consists of an optimal opponent.

In a game, there are two or more independently-acting players who make choices (moves) and receive a payoff based on the choices they make. The differential "game of simple pursuit" is a two-person, competitive, zero-sum, simultaneous-moving, complete-information game in which a fast pursuing player $P$ is trying to capture a slower evading player $E$. The "choices" available to a player at a given moment consist of choosing a direction (angle) in which to travel. In the simple game of pursuit, the players travel in a plane and both players may instantaneously change direction. Each player travels at a constant speed, and the pursuing player’s speed $w_P$ (1.0) is greater than the evading.
player’s speed $w_e (0.67)$.

The state variables of the game are $x_P, y_P, x_e, y_e$ representing the coordinate positions $(x_P, y_P)$ and $(x_e, y_e)$ of the pursuer $P$ and evader $E$ in the plane.

At each time step, both players know the position (state variables) of both players. The choice for each player is to select a value of their control variable (i.e. the angular direction in which to travel). The pursuer’s control variable is the angle $\phi$ (from 0 to $2\pi$ radians) and the evader’s control variable is the angle $\psi$. The players choose their respective control variable simultaneously.

The analysis of this game can be simplified by reducing the number of state variables from four to two (Isaacs 1965). This state reduction is accomplished by simply viewing the pursuer $P$ as being at the origin point $(0,0)$ of a new coordinate system at all times and then viewing the evader $E$ as being at position $(x, y)$ in this new coordinate system. The two numbers $x$ and $y$ representing the position $(x, y)$ of the evader $E$ thus become the two “reduced” state variables of the game. Whenever the pursuer $P$ travels in a particular direction, the coordinate system is immediately adjusted so that the pursuer is repositioned back to the origin $(0,0)$ and then appropriately adjusting the position $(x, y)$ of the evader to reflect the travel of the pursuer.

The state transition equations for the evader $E$ follow:

$$x(t+1) = x(t) + w_e \cos \psi - w_p \cos \phi$$

$$y(t+1) = y(t) + w_e \sin \psi - w_p \cos \phi.$$

In order to develop optimal playing strategies, we use a set of random environmental starting condition cases consisting of $N_e (= 10)$ starting positions $(x_i, y_i)$ for the evader $E$. Each starting value of $x_i$ and $y_i$ is between -5.0 and +5.0. The two players may travel anywhere in the plane. We regard the pursuer $P$ as having captured the evader $E$ when the pursuer gets to within some specified small distance $\varepsilon = 0.5$ of the evader $E$.

The payoff for a given player is measured by time. The payoff for the pursuer $P$ is the total time it takes to capture the evader $E$ over all of the environmental cases. The pursuer tries to minimize the time to capture. The payoff for the evader $E$ is the total time of survival for $E$. The evader tries to maximize this time of survival.

A maximum “time out” time (100 time steps) is established so that if a particular pursuer strategy has not made the capture within that amount of time, that maximum time becomes the payoff for that particular environmental case and that particular strategy.

The problem is to find the strategy for choosing the control variable of the pursuer so as to minimize the total time to capture for any set of environmental cases when playing against an optimal evader.

For this particular simple game, the best strategy for the pursuer $P$ at any given time step is to chase the evader $E$ in the direction of the straight line currently connecting the pursuer to the evader. And, for this particular simple game, the best strategy for the evader $E$ is to race away from the pursuer in the direction of the straight line connecting the pursuer to the evader.

If the evader chooses some action other than the strategy of racing away from the pursuer in the direction of the straight line connecting the pursuer to the evader (as shown below), the evader will survive for less time than if he follows his best strategy. If the evader initially chooses a sub-optimal direction and then belatedly chooses the optimal direction, his time of survival is less than if he had chosen the optimal direction from the beginning.

The situation is symmetric in that if the pursuer does not chase after the evader $E$ along the straight line, he fails to minimize the time to capture.

The “value of the game” is the payoff (time) such that, no matter what the evader does, the evader cannot hold out for longer than this amount of time, and, if the evader does anything other than direct fleeing, his survival time is a shorter amount of time. Conversely, no matter what the pursuer does, the pursuer $P$ cannot capture an optimal evader $E$ in less than that amount of time. And, if the pursuer does anything other than direct pursuit, the evader can remain at large for a longer amount of time.

The genetic programming paradigm can be used to solve
the differential game of simple pursuit by genetically evolv-
ing a population of strategies for the pursuing individuals
over a number of generations.

The genetic programming paradigm is especially well
suited to solving this kind of problem because the solution
takes the form of a mathematical expression whose size and
shape may not be known in advance.

The terminal set is \( T = \{X, Y, R\} \). The two state vari-
ables \( X \) and \( Y \) represent the position of the evader \( E \) in the
plane in a "reduced" coordinate system where the pursuer is
always positioned (or repositioned) to the origin. The ter-

inal \( R \) is an ephemeral random constant. Each occurrence
of an ephemeral random constant in an individual in the ini-
tial random population is assigned a different random float-
ing point value (between -1.0 and +1.0). These ephemeral
random constants give the algorithm the building blocks
with which to construct the particular numeric constants that
may be needed in the solution. Additional details are found
in Koza (1990a).

The function set for this problem can be a set of arith-
metic and mathematical operations such as addition, subtrac-
tion, multiplication, division (using the operation \( \% \) which
returns a zero when division by zero is attempted), and the
exponential function \( \text{EXP} \). Thus, the function set is \( F = \{+, \,-, \,* , \,\% , \,\text{EXP}\} \).

If the population of individuals represents pursuers and
we are attempting to genetically breed an optimal pursuing
individual, the environment for this "genetic algorithm"
consists of an optimal evading individual. The optimal
evader travels with the established constant evader speed \( w_e \)
in the angular direction specified by the Arctangent function.

The genetic programming paradigm is successful in ge-
etically breeding optimal pursuing individuals. As one pro-
gresses from generation to generation, the population of
pursuing individuals typically improves. After several gener-
ations, the best pursuing individuals in the population can
capture the evader in a small fraction (perhaps 2, 3, or 4) of
the 10 environmental cases within a certain amount of time.
Then, after additional generations, the population improves
and the best pursuing individuals in the population can cap-
ture the evader in a larger fraction (perhaps 4, 5, or 6) of the
10 environmental cases within a shorter amount of time.
Often, these partially effective pursuers are effective in some
identifiable fraction of the plane or at some identifiable range
of distances, but ineffective in other parts of the plane or at
other distances. However, as more and more generations are
run, the population of pursuing individuals typically contin-
ues to improve.

In one run, a pursuer strategy emerged in the 17th gen-
eration which correctly handled all 10 of the environmental
cases. This S-expression is shown below.

\[
(\% (- (\% (+ (* 2.0 Y) -0.066) -0.365) (% Y -0.124)) (+ (\text{EXP} X) Y -0.579))
\]

An optimal evader has been similarly evolved using an
optimal pursuer (i.e. the Arctangent strategy).

6. CO-EVOLUTION IN NATURE

The evolutionary process in nature is often described as if
one population of individuals is trying to adapt to a fixed
environment. This description is, however, only a first order
approximation to the actual situation. The environment ac-
tually consists of both the physical environment (which is usually relatively unchanging) as well as other independently-acting biological populations of individuals which are simultaneously trying to adapt to “their” environment. The actions of each of these other independently-acting biological populations (species) usually affect all the others. In other words, the environment of a given species includes all the other biological species that contemporaneously occupy the physical environment and which are simultaneously trying to survive. In biology, the term “co-evolution” is sometimes used to reflect the fact that all species are simultaneously co-evolving in a given physical environment.

A biological example presented by Holland illustrates the point (1990). A given species of plant may be faced with an environment containing insects that like to eat it. To defend against its predators (and increase its probability of survival in the environment), the plant may, over a period of time, evolve a tough exterior that makes it difficult for the insect to eat it. But, over a period of time, the insect may evolve a stronger jaw so that the insect population can continue to feed on the plant (and increase its probability of survival in the environment). Then, over an additional period of time, the plant may evolve a poison to help defend itself further against the insects. But, then again, over a period of time, the insect may evolve a digestive enzyme that negates the effect of the poison so that the insect population can continue to feed on the plant.

In effect, both the plant and the insects get better and better at their respective defensive and offensive roles in this “biological arms race”. Each species changes in response to the actions of the other.

7. BACKGROUND ON CO-EVOLUTION AND GENETIC ALGORITHMS

In the “genetic algorithm,” described by John Holland in his pioneering Adaptation in Natural and Artificial Systems (1975), a population of individuals attempts to adapt to a fixed “environment.” In the basic genetic algorithm as described by Holland in 1975, the individuals in the population are fixed-length character strings (typically binary strings) that are encoded to represent some problem in some way. In the basic “genetic algorithm”, the performance of the individuals in the population is measured using a fitness measure which is, in effect, the “environment” for the population. Over a period of many generations, the genetic algorithm causes the individuals in the population to adapt in a direction that is dictated by the fitness measure (its environment).

Holland (1990) has incorporated co-evolution and genetic algorithms in his ECHO system for exploring the co-evolution of artificial organisms described by fixed-length character strings (chromosomes) in a “miniature world.” In ECHO, there is a single population of artificial organisms. The environment of each organism includes all other organisms.

Miller (1988, 1989) has used co-evolution in a genetic algorithm to evolve a finite automaton as the strategy for playing the Repeated Prisoner’s Dilemma game. Miller’s population consisted of strings (chromosomes) of 148 binary digits to represent finite automata with 16 states. Each string in the population represented a complete strategy by which to play the game. That is, it specified what move the player was to make for any sequence of moves by the other player. Miller then used co-evolution to evolve strategies. Miller’s co-evolutionary approach contrasts with Axelrod’s (1987) solution to the repeated prisoner’s dilemma using genetic algorithms. Axelrod measured performance of a particular strategy with a fixed weighted mix of the strategy’s results against eight superior computer programs submitted in an international programming tournament for the prisoner’s dilemma. A best strategy for one player (represented as a 70 bit string with a 3-move look-back) was then evolved with the weighted mix of eight opposing computer programs serving as the environment. Hillis (1990) used co-evolution in genetic algorithms to solve optimization problems.

8. CO-EVOLUTION AND THE GENETIC PROGRAMMING PARADIGM

In the “hierarchical co-evolution algorithm,” there are two populations of individuals. The environment for the first population consists of the second population. And, conversely, the environment for the second population consists of the first population.

The co-evolutionary process typically starts with both populations being highly unfit (when measured by an absolute fitness measure). Then, the first population tries to adapt to the “environment” created by the second population. Simultaneously, the second population tries to adapt to the “environment” created by the first population.

This process is carried out by testing the performance of each individual in the first population against each individual (or a sampling of individuals) from the second population. We call this performance the “relative fitness” of an individual because it represents the performance of one individual in one population relative to the environment consisting of the entire second population. Then, each individual in the second population is tested against each individual (or a sampling of individuals) from the first population.

Note that this measurement of relative fitness for an individual in co-evolution is not an absolute measure of fitness against an optimal opponent, but merely relative measure when the individual is tested against the current opposing population. If one population contains boxers who only throw left punches, then an individual whose defensive repertoire contains only defenses against left punches will have high relative fitness. But, this individual will have only mediocre absolute fitness when tested against an opponent who knows how to throw both left punches and right punches (the optimal opponent).

Even though both initial populations are initially highly unfit (both relatively and absolutely), the virtually inevitable variation of an initial random population will mean that some individuals have slightly better relative fitness than others. That means that some individuals in each population have somewhat better performance than others in dealing with the current opposing population.

The operation of fitness proportionate reproduction (based on the Darwinian principle of “survival and reproduc-
tion of the fittest) can then be applied to each population using the relative fitness of each individual currently in each population. In addition, the operation of genetic recombination (crossover) can also be applied to a pair of parents, at least one of which is selected based on its relative fitness.

Over a period of time, both populations of individuals will tend to “co-evolve” and to rise to higher levels of performance as measured in terms of absolute fitness. Both populations do this without the aid of any externally supplied absolute fitness measure serving as the environment. In the limiting case, both populations of individuals can evolve to a level of performance that equals the absolute optimal fitness. Thus, the hierarchical co-evolution algorithm is a self-organizing, mutually-bootstrapping process that is driven only by relative fitness (and not by absolute fitness).

Co-evolution is likely to be especially important in game theory because one almost never has a priori access to a minimax strategy for either player. One therefore encounters a "chicken and egg" situation. In trying to develop a minimax strategy for the first player, one does not have the advantage of having a minimax second player against which to test candidate strategies. In checkers or chess, for example, it is difficult for a new player to learn to play well if he does not have the advantage of playing against a highly competent player.

9. CO-EVOLUTION OF A GAME STRATEGY

We now illustrate the “hierarchical co-evolution algorithm” to discover minimax strategies for both players simultaneously in a simple discrete game represented by a game tree in extensive form.

In the hierarchical co-evolution algorithm, we do not have access to the optimal opponent to train the population. Instead, our objective is to breed two populations simultaneously. Both populations start as random compositions of the available functions and arguments. The entire second population serves as the “environment” for testing the performance of each particular individual in the first population. And, at the same time, the entire first population serves as the environment for testing the performance of each particular individual in the second population. The algorithm does not have access to the “absolute fitness” measure provided by an optimal opponent. In other words, only relative fitness is used to drive the co-evolution algorithm.

Consider the following simple discrete game whose game tree is presented in extensive form in Figure 2. Each internal point of this tree is labeled with the player who must move. Each line is labeled with the choice (either L or R) made by the moving player. Each endpoint of the tree is labeled with the payoff (to player X).

This game is a two-person, competitive, zero-sum game in which the players make alternating moves. On each move, a player can choose to go L (left) or R (right). After player X has made three moves and player O has made two moves, player X receives (and player O pays out) the payoff shown at the particular endpoint of the game tree (1 of 32).

Since this game is a game of complete information, each player has access to complete information about his opponent's previous moves (and his own previous moves). This historical information is contained in five variables XM1 (X’s move 1), OM1 (O’s move 1), XM2 (X’s move 2), OM2 (O’s move 2), and XM3 (X’s move 3). These five variables each assume one of three possible values: L (left), R (right), or U (undefined). A variable is undefined prior to the time when the move to which it refers has been made. Thus, at the beginning of the game, all five variables are undefined. The particular variables that are defined and undefined indicate the point to which play has progressed during the play of the game. For example, if both players have moved once, XM1 and OM1 are defined (as either L or R) but the other three variables (XM2, OM2, and XM3) are undefined (have the value U).

A strategy for a particular player in a game specifies which choice that player is to make for every possible situation that may arise for that player. In particular, in this game, a strategy for player X must specify his first move if he happens to be at the beginning of the game. A strategy for player X must also specify his second move if player O

![Figure 2 Game Tree with Payoffs](image-url)
has already made one move and it must specify his third move if player O has already made two moves. Since Player X moves first, player X's first move is not conditioned on any previous move. But, player X's second move will depend on Player O's first move (i.e. OM1) and, in general, it will also depend on his own first move (XM1). Similarly, player X's third move will depend on player O's first two moves and, in general, his own first two moves. Similarly, a strategy for player O must specify what choice player O is to make for every possible situation that may arise for player O. A strategy here is a computer program (i.e. S-expression) whose inputs are the relevant historical variables and whose output is a move (L or R) for the player involved. Thus, the set of terminals is T = {L, R}.

Four testing functions CXM1, COM1, CXM2, and COM2 provide the facility to test each of the historical variables that are relevant to deciding upon a player's move. Each of these functions is a specialized form of the CASE function in LISP. For example, function CXM1 has three arguments and evaluates its first argument if XM1 (X’s move 1) is undefined, evaluates its second argument if XM1 is L (Left), and evaluates its third argument if XM1 is R (Right). Functions CXM2, COM1, and COM2 are similarly defined. Thus, the function set for this problem is F = {CXM1, COM1, CXM2, COM2}. Each of these functions takes three arguments.

Our goal is to simultaneously co-evolve strategies for both players of this game.

In co-evolution, the relative fitness of a particular strategy for a particular player in a game is the average of the payoffs received when that strategy is played against the entire population of opposing strategies.

The absolute fitness of a particular strategy for a particular player in a game is the payoff received when that strategy is played against the minimax strategy for the opponent. Note that when we compute the absolute fitness of an X strategy for our descriptive purposes here, we test the X strategy against 4 possible combinations of O moves — that is, O’s choice of L or R for his moves 1 and 2. When we compute the absolute fitness of an O strategy, we test it against 8 possible combinations of X moves — that is, X’s choice of L or R for his moves 1, 2, and 3. Note that this testing of 4 or 8 combinations does not occur in the computation for relative fitness. When the two minimax strategies are played against each other, the payoff is 12, which is the value of this game. A minimax strategy takes advantage of non-minimax play by the other player.

As previously mentioned, the co-evolution algorithm does not use the minimax strategy of the opponent in any way. We use it in this paper for descriptive purposes. The co-evolution algorithm uses only relative fitness.

In one run (with population size 300), the individual strategy for player X in the initial random generation (generation 0) with the best relative fitness was

\[(\text{COM1 L} \ (\text{COM2 L L R}) \ R) \ .\]

This individual has relative fitness of 10.08.

The individual in the initial random population (generation 0) for player O with the best relative fitness was an equally complex expression. It simplifies to

\[\text{COM2 R} \ (\text{COM1 # L R}) \ (\text{COM1 # R L})] \ .\]

Note that, in simplifying this strategy, we inserted the symbol # to indicate that the situation involved can never arise. This individual has relative fitness of 7.57.

Neither the best X individual nor the best O individual from generation 0 reached maximal absolute fitness.

Note that the values of relative fitness for the relative best X individual and the relative best O individual from generation 0 (i.e. 10.08 and the 7.57) are each computed by averaging the payoff from the interaction of the individual involved with all 300 individual strategies in the current opposing population.

In generation 1, the individual strategy for player X with the best relative fitness had relative fitness of 11.28. This individual X strategy is still not a minimax strategy. It does not have the maximal absolute fitness.

In generation 1, the best individual O strategy attained relative fitness of 7.18. It is shown below:

\[\text{COM2} \ (\text{COM1 R R L}) \ (\text{COM2 L L (COM2 R L R)}) \ R] \ .\]

This individual O strategy simplifies to

\[\text{COM2} \ (\text{COM1 R R L}) \ L R] \ .\]

Although the co-evolution algorithm does not know it, this best single individual O strategy for generation 1 is a minimax strategy for player O. It has maximal absolute fitness in this game. This one O individual was the first such O individual to attain this level of performance during this run. If it were played against the minimax X strategy, it would score 12 (i.e. the value of this game).

Between generations 2 and 14, the number of individuals in the O population reaching maximal absolute fitness was 2, 7, 17, 28, 35, 40, 50, 64, 73, 83, 93, 98, and 107, respectively. That is, the minimax O strategy began to dominate the O population.

In generation 14, the individual strategy for player X with the best relative fitness had relative fitness of 18.11. Although the co-evolution algorithm does not know it, this one individual scored maximal absolute fitness. This one X individual was the first such X individual to attain this level of performance during this run. If it were played against the minimax O strategy, it would score 12 (i.e. the value of this game).

Between generations 2 and 14, the number of individuals in the O population reaching maximal absolute fitness was 2, 7, 17, 28, 35, 40, 50, 64, 73, 83, 93, 98, and 107, respectively. That is, the minimax O strategy began to dominate the O population.

This individual X strategy simplifies to

\[\text{COM2} \ (\text{COM1 L L R}) \ L R] \ .\]

Between generations 15 and 29, the number of individuals in the X population reaching maximal absolute fitness was 3, 4, 8, 11, 10, 9, 13, 21, 24, 29, 43, 32, 52, 48, and 50, respectively. That is, the minimax X strategy began to dominate the X population. Meanwhile, the O population became even more dominated by the O minimax strategy.

By generation 38, the number of O individuals in the population reaching maximal absolute fitness reached 188.
(almost two thirds of the population) and the number of X individuals reaching maximal absolute fitness reached 74 (about a quarter). That is, by generation 38, the minimax strategies for both players were dominant.

Interestingly, these 74 X individuals had relative fitness of 19.11 and these 188 O individuals had relative fitness of 10.47. Neither of these values equals 12 because the other population is not fully converged to its minimax strategy.

In summary, we have seen the discovery, via mutual bootstrapping, of the minimax strategies for both players in this game. This mutually bootstrapping process found the minimax strategies for both players without using knowledge of the minimax strategy (i.e. any a priori knowledge of the game) for either player.

10. CONCLUSIONS

We used the genetic programming paradigm to discover an "artificial ant" to traverse an irregular trail and to discover the minimax strategy for a pursuer in the differential game of simple pursuit. We then extended the genetic programming paradigm so that two populations of individuals are simultaneously co-evolved wherein each population serves as the environment to guide the evolution of the other population. We then illustrated the hierarchical co-evolution algorithm by simultaneously breeding an optimal game-playing strategy for both players of a discrete game in extensive form.

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REFERENCES


